

Complete ZX-calculi for the stabiliser fragment in odd prime dimensions

Robert I. Booth, Titouan Carette

QPL 2022 (...also MFCS 2022)



THE UNIVERSITY of EDINBURGH
informatics

- State space $\rightarrow p$ -dimensional Hilbert space:

$$\mathcal{H} = L^2(\mathbb{Z}_p) \cong \mathbb{C}^d$$

$$|\psi\rangle \in \mathcal{H} \iff |\psi\rangle = \sum_{x \in \mathbb{Z}_p} c_x |x\rangle$$

- p is prime $\implies \mathbb{Z}_p$ is a field

Quantum computing with qubits: gates

- Pauli gates

$$Z|x\rangle = e^{\frac{i2\pi x}{p}}|x\rangle = \omega^x|x\rangle \quad X|x\rangle = |x+1 \bmod p\rangle$$

$$\text{generic Pauli} \rightsquigarrow X^a Z^b \quad \text{for } a, b \in \mathbb{Z}_p$$

Quantum computing with qubits: gates

- Pauli gates

$$Z|x\rangle = e^{\frac{i2\pi x}{p}}|x\rangle = \omega^x|x\rangle \quad X|x\rangle = |x+1 \bmod p\rangle$$

$$\text{generic Pauli} \rightsquigarrow X^a Z^b \quad \text{for } a, b \in \mathbb{Z}_p$$

- Hadamard gate

$$H|x\rangle = \sum_{k \in \mathbb{Z}_p} \omega^{kx} |x\rangle \quad \text{and} \quad HZH^\dagger = X$$

Quantum computing with qubits: gates

- Pauli gates

$$Z|x\rangle = e^{\frac{i2\pi x}{p}}|x\rangle = \omega^x|x\rangle \quad X|x\rangle = |x+1 \bmod p\rangle$$

$$\text{generic Pauli} \rightsquigarrow X^a Z^b \quad \text{for } a, b \in \mathbb{Z}_p$$

- Hadamard gate

$$H|x\rangle = \sum_{k \in \mathbb{Z}_p} \omega^{kx} |k\rangle \quad \text{and} \quad HZH^\dagger = X$$

- Controlled phase or CZ-gate

$$E|x\rangle \otimes |y\rangle = \omega^{xy} |x\rangle \otimes |y\rangle$$

Quantum computing with qubits: gates

- Pauli gates

$$Z|x\rangle = e^{\frac{i2\pi x}{p}}|x\rangle = \omega^x|x\rangle \quad X|x\rangle = |x+1 \bmod p\rangle$$

$$\text{generic Pauli} \rightsquigarrow X^a Z^b \quad \text{for } a, b \in \mathbb{Z}_p$$

- Hadamard gate

$$H|x\rangle = \sum_{k \in \mathbb{Z}_p} \omega^{kx} |k\rangle \quad \text{and} \quad HZH^\dagger = X$$

- Controlled phase or CZ-gate

$$E|x\rangle \otimes |y\rangle = \omega^{xy} |x\rangle \otimes |y\rangle$$

- Phase gate

$$S|x\rangle = \omega^{2^{-1}x(x+1)} |x\rangle$$

Quantum computing with qubits: gates (cont.)

These gates are **not** self-inverse

$$(X^a Z^b)^p = 1_{\mathcal{H}} \quad S^p = 1_{\mathcal{H}} \quad E^p = 1_{\mathcal{H} \otimes \mathcal{H}} \quad H^4 = 1_{\mathcal{H}}.$$

The stabiliser fragment

- Recall the qubit stabiliser fragment:

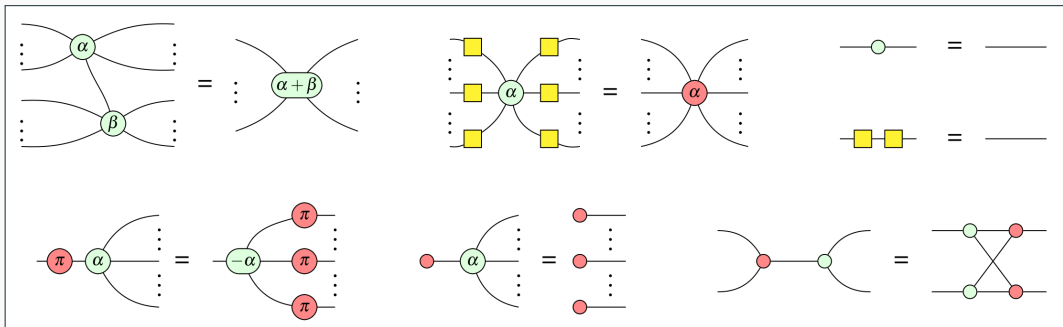
Pauli states + Clifford unitaries + Pauli projections

- The Clifford group on n qubits is the normaliser in $U(\mathcal{H}^{\otimes n})$ of the group of Paulis.

- Formally, the Clifford group is a representation of $\text{ASymp}(\mathbb{Z}_p^{2n}) \rtimes \mathbb{T}$

Symplectic matrix + Translation + Global phase

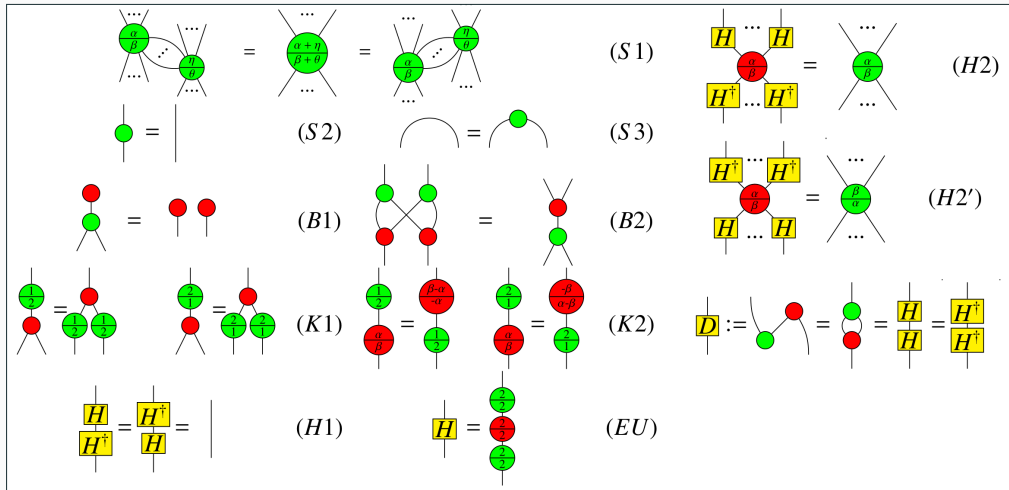
Qubit stabiliser ZX



Miriam Backens. "The ZX-calculus Is Complete for Stabilizer Quantum Mechanics". In: *New Journal of Physics* 16.9 (Sept. 17, 2014), p. 093021.

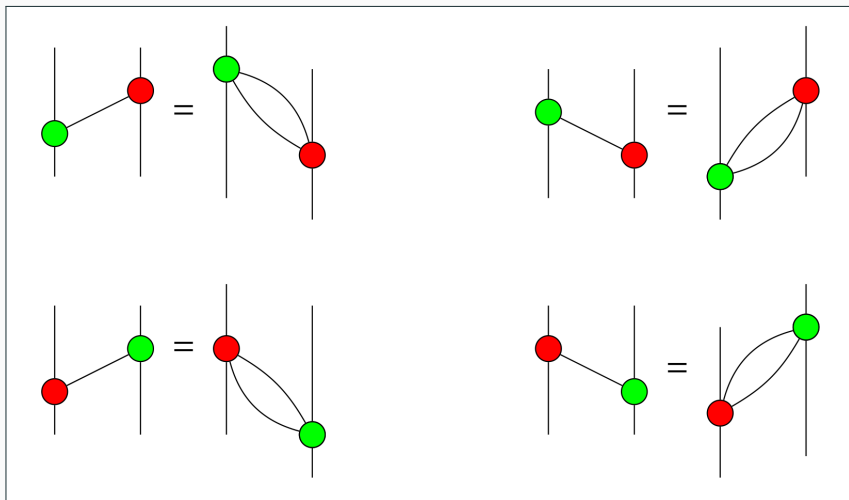
Tommy McElvanney and Miriam Backens. *Complete Flow-Preserving Rewrite Rules for MBQC Patterns with Pauli Measurements*. May 4, 2022.

Qutrit stabiliser ZX



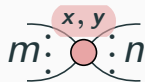
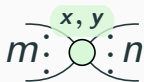
Quanlong Wang. "Completeness of the ZX-calculus". PhD thesis. Oxford: Oxford University, 2018.

Qutrit stabiliser ZX



Quanlong Wang. "Completeness of the ZX-calculus". PhD thesis. Oxford: Oxford University, 2018.

Generators of the calculus



...for $m, n \in \mathbb{N}$ and $x, y \in \mathbb{Z}_p$.

Sequential and parallel composition defined in the usual way.

Standard interpretation

$$\left[\begin{array}{c} \text{green circle with } x, y \text{ above it} \\ m \vdots n \end{array} \right] = \sum_{k \in \mathbb{Z}_p} \omega^{2^{-1}(xk+yk^2)} |k : Z\rangle^{\otimes n} \langle k : Z|^{\otimes m}$$

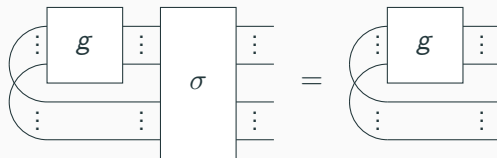
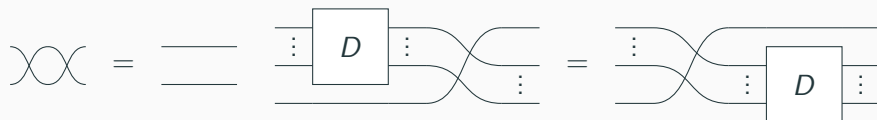
$$\left[\begin{array}{c} \text{red circle with } x, y \text{ above it} \\ m \vdots n \end{array} \right] = \sum_{k \in \mathbb{Z}_p} \omega^{2^{-1}(xk+yk^2)} |-k : X\rangle^{\otimes n} \langle k : X|^{\otimes m}$$

$$\left[\text{yellow square} \right] = H$$

$$\left[\begin{array}{c} \text{green circle with } 1, 1 \text{ above it} \\ \text{green circle} \end{array} \right] = S$$

$$\left[\begin{array}{c} \text{red circle} \\ \text{green circle} \\ \text{yellow square} \\ \text{green circle} \end{array} \right] = E$$

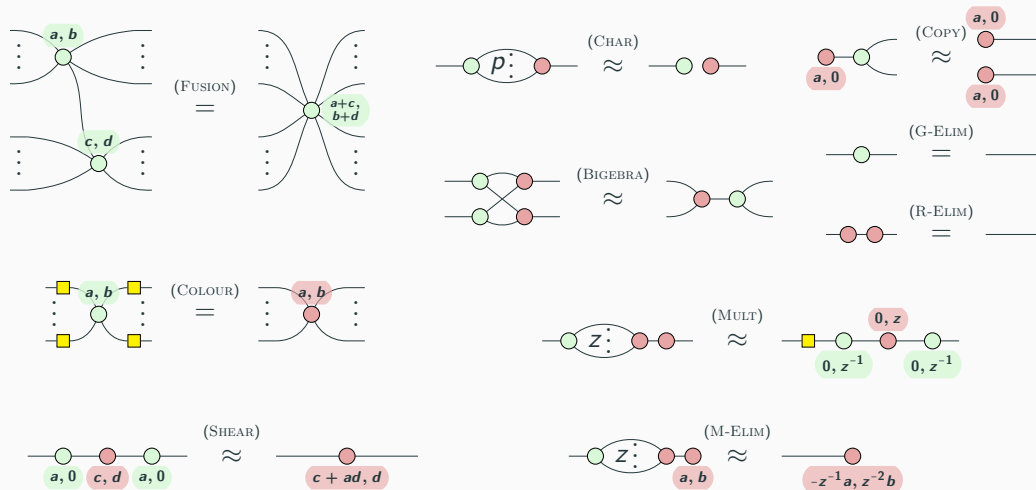
Structural axioms



The calculus is flexsymmetric!

Titouan Carette. "Wielding the ZX-calculus, Flexsymmetry, Mixed States, and Scalable Notations". Theses. Université de Lorraine, Nov. 2021.

Scalarless axiomatisation



This calculus is **sound**, **universal** and **complete** for the stabiliser fragment in any odd prime dimension.

Things I don't have time to talk about

- Meta-rule for “spider wars”

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers
- **Weighted Hadamards: flexsymmetrisation of multipliers**

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers
- Weighted Hadamards: flexsymmetrisation of multipliers
- A natural representation of graph states

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers
- Weighted Hadamards: flexsymmetrisation of multipliers
- A natural representation of graph states
 - Pauli stabilisers

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers
- Weighted Hadamards: flexsymmetrisation of multipliers
- A natural representation of graph states
 - Pauli stabilisers
 - Local complementation

Things I don't have time to talk about

- Meta-rule for “spider wars”
- A version of the colour change rule
- Syntactic sugar for multi-edges: multipliers
- Weighted Hadamards: flexsymmetrisation of multipliers
- A natural representation of graph states
 - Pauli stabilisers
 - Local complementation
- Completeness of the scalar fragment

Mixed states and the CPM construction

- Straightforward extension to a complete axiomatisation for mixed stabiliser states using the **discard construction**.
- An alternative interpretation in terms of affine co-isotropic relations over $\mathbb{Z}_p$, which is sound and complete:

$$\begin{array}{ccc} ZX_p & \xrightarrow{\quad [-] \quad} & \text{AffColsoRel}_{\mathbb{Z}_p} \\ \begin{array}{c} \downarrow \\ \llbracket - \rrbracket \end{array} & & \begin{array}{c} \updownarrow \end{array} \\ \text{CPM}(\text{Stab}_p) & \longrightarrow & \text{CPM}(\text{Stab}_p)/(\text{SCALARS}) \end{array}$$

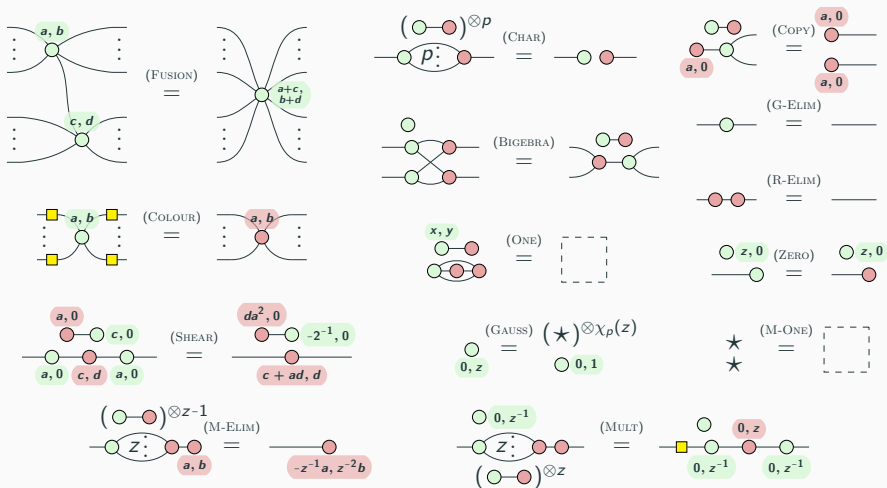
Titouan Carette et al. “Completeness of Graphical Languages for Mixed State Quantum Mechanics”. In: *ACM Transactions on Quantum Computing* 2.4 (Dec. 21, 2021), 17:1–17:28.

Cole Comfort and Aleks Kissinger. “A Graphical Calculus for Lagrangian Relations”. May 13, 2021.

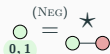
- Cleaner/simplified axiomatisation
- Complete ZX-calculi for all of quantum theory in prime dimensions?
- Complete ZX-calculi for stabiliser theory in non-prime dimensions?
- What about affine co-isotropic relations over other fields?

Thanks for listening!

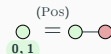
Full axiomatisation



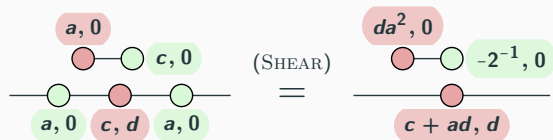
if $p \equiv 1 \pmod{8}$



if $p \equiv 5 \pmod{8}$

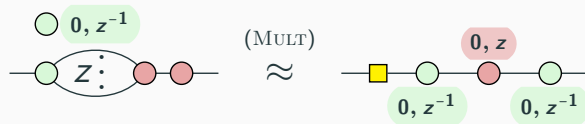
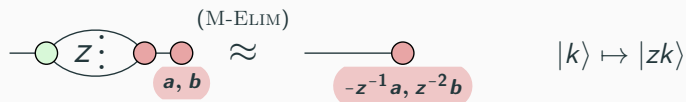


Axiom for affine shears



$$\begin{bmatrix} 1 & 0 \\ d & 1 \end{bmatrix} \left(\begin{bmatrix} x \\ y \end{bmatrix} + \begin{bmatrix} 0 \\ a \end{bmatrix} \right) = \begin{bmatrix} 1 & 0 \\ d & 1 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} + \begin{bmatrix} ad \\ a \end{bmatrix}$$

Axioms for multipliers



$$\begin{bmatrix} z & 0 \\ 0 & z^{-1} \end{bmatrix} = \begin{bmatrix} 1 & z^{-1} \\ 0 & 1 \end{bmatrix} \begin{bmatrix} -1 & 0 \\ -z & -1 \end{bmatrix} \begin{bmatrix} 1 & z^{-1} \\ 0 & 1 \end{bmatrix} \begin{bmatrix} 0 & -1 \\ 1 & 0 \end{bmatrix}$$

Scalar axioms

$$\begin{array}{c} x, y \\ \text{---} \circ \text{---} \circ \\ \text{---} \circ \text{---} \circ \end{array} \stackrel{(\text{ONE})}{=} \boxed{\phantom{\text{---} \circ \text{---} \circ}}$$

$$\begin{array}{c} \circ \text{---} \circ \\ \text{---} \circ \end{array} \stackrel{(\text{ZERO})}{=} \begin{array}{c} \circ \text{---} \circ \\ \text{---} \circ \end{array}$$

$$\begin{array}{c} \circ \\ \text{---} \circ \end{array} \stackrel{(\text{GAUSS})}{=} \begin{array}{c} \circ \\ \text{---} \circ \end{array} \otimes \chi_p(z)$$

$$\begin{array}{c} \star \\ \star \end{array} \stackrel{(\text{M-ONE})}{=} \boxed{\phantom{\text{---} \circ \text{---} \circ}}$$

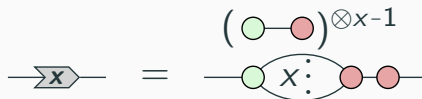
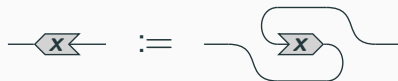
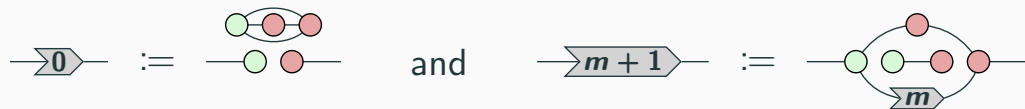
if $p \equiv 1 \pmod{8}$

$$\begin{array}{c} \circ \\ \text{---} \circ \end{array} \stackrel{(\text{NEG})}{=} \begin{array}{c} \star \\ \text{---} \circ \end{array}$$

if $p \equiv 5 \pmod{8}$

$$\begin{array}{c} \circ \\ \text{---} \circ \end{array} \stackrel{(\text{Pos})}{=} \begin{array}{c} \circ \text{---} \circ \end{array}$$

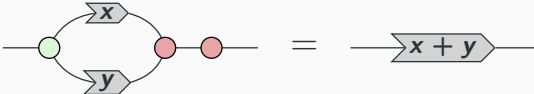
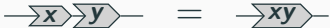
Syntactic sugar for multi-edges



Multipliers are **not** flexsymmetric!

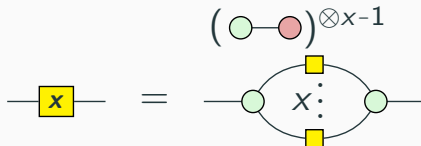
Fabio Zanasi. "Interacting Hopf Algebras: The Theory of Linear Systems". PhD thesis. Ecole Normale Supérieure de Lyon, May 4, 2018.

Syntactic sugar for multi-edges (cont.)



Multipliers are a presentation of the field $\mathbb{Z}_p$!

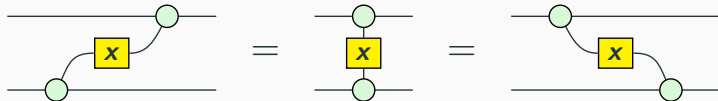
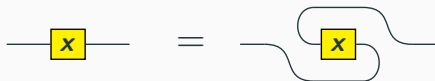
Weighted Hadamards



Alex Townsend-Teague and Konstantinos Meichanetzidis. "Classifying Complexity with the ZX-Calculus: Jones Polynomials and Potts Partition Functions". Mar. 11, 2021.

Weighted Hadamards (cont.)

Weighted Hadamards **are** flexsymmetric:



Weighted Hadamard equations

$$\text{---} \boxed{x} \boxed{y} \text{---} = \text{---} \boxed{-xy^{-1}} \text{---}$$

$$\text{---} \boxed{x} \boxed{-x} \text{---} = \text{---}$$

$$\text{---} \boxed{0} \text{---} = \text{---} \bigcirc \bigcirc \text{---}$$

$$\text{---} \bigcirc \begin{array}{c} \boxed{x} \\ \boxed{y} \end{array} \bigcirc \text{---} = \text{---} \boxed{x+y} \text{---}$$

$$\text{---} \boxed{x} \boxed{x} \boxed{x} \text{---} = \text{---} \boxed{-x} \text{---}$$

$$\text{---} \boxed{1} \text{---} = \text{---} \boxed{} \text{---}$$